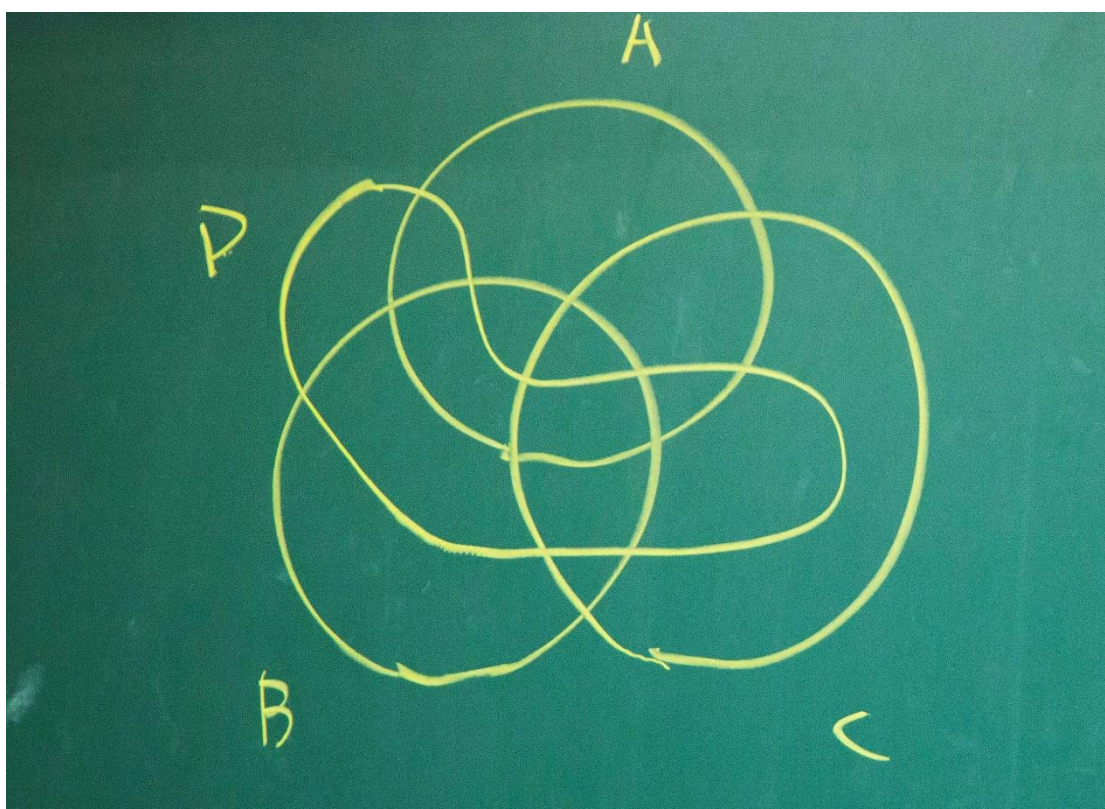
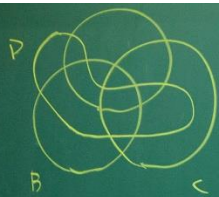


【10920 趙啟超教授離散數學 / 第 3 堂版書】





Mathematical Induction

induction step

induction basis

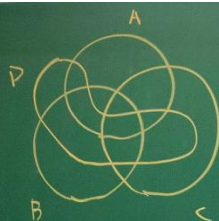
induction hypothesis

Principle of mathematical induction:

Suppose that $S(n)$ is a statement involving the variable n with the following properties:

(i) $S(1)$ is true;

(ii) if $S(k)$ is true, then $S(k+1)$ is true, for every $k \in \mathbb{N}$.



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Example Show that for all $n \in \mathbb{N}$,

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Proof Induction basis: When $n=1$,

$$\sum_{\lambda=1}^1 (4\lambda+3) = 7 = 2 \cdot 1^2 + 5 \cdot 1$$

Induction step: Assume that $\sum_{\lambda=1}^k (4\lambda+3) = 2k^2+5k$.

$$\begin{aligned} \text{Then } \sum_{\lambda=1}^{k+1} (4\lambda+3) &= \sum_{\lambda=1}^k (4\lambda+3) + 4(k+1)+3 \\ &= 2k^2+5k + (4k+7) \end{aligned}$$

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$$= 2k^2 + 9k + 7$$

$$= 2(k^2 + 2k + 1) + 5k + 5$$

$$= 2(k+1)^2 + 5(k+1)$$

By mathematical induction, this holds for all $n \in \mathbb{N}$.

Example Show that for $n \in \mathbb{N}$.

$$\sum_{i=1}^n i^3 = \left(\frac{n(n+1)}{2} \right)^2$$

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$$\begin{aligned}
 \text{Then } \sum_{i=1}^{k+1} i^3 &= \sum_{i=1}^k i^3 + (k+1)^3 \\
 &= \left(\frac{k(k+1)}{2}\right)^2 + (k+1)^3 \\
 &= (k+1)^2 \left(\frac{k^2}{4} + k + 1\right) \\
 &= (k+1)^2 \frac{1}{4} (k^2 + 4k + 4) \\
 &= \left(\frac{(k+1)(k+2)}{2}\right)^2
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By mathematical induction, $\sum_{i=1}^n i^3 = \left(\frac{n(n+1)}{2}\right)^2$ for all $n \in \mathbb{N}$. \square

Let $a_n = \sum_{i=1}^n i^3$. Then $a_n - a_{n-1} = n^3$ with $a_1 = 1$.
for $n \geq 2$.

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Principle of mathematical induction - Alternative form

(i) $S(n_0), S(n_0+1), \dots, S(n_1)$ are true;

(ii) if $S(n_0), S(n_0+1), \dots, S(k)$ are true for $k \geq n_1$,

then $S(k+1)$ is also true.

Then $S(n)$ is true for all $n \geq n_0$.

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if defined by $a_1 = 1, a_2 = 5$
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Then $P(n)$ is true for all $n \geq n_0$.

Example Show that if a_n is defined by $a_1 = 1, a_2 = 5$,
and $a_{n+1} = 5a_n - 6a_{n-1}$ for $n \geq 2$,
then $a_n = 3^n - 2^n$ for all $n \in \mathbb{N}$.

Proof Induction basis: For $n=1, 3^1 - 2^1 = 1 = a_1$.
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$$\begin{aligned} \text{Then } a_{k+1} &= 5a_k - 6a_{k-1} \\ &= 5(3^k - 2^k) - 6(3^{k-1} - 2^{k-1}) \\ &= (5 \cdot 3 - 6)3^{k-1} - (5 \cdot 2 - 6)2^{k-1} \\ &= 9 \cdot 3^{k-1} - 4 \cdot 2^{k-1} = 3^{k+1} - 2^{k+1}. \end{aligned}$$

Therefore, by mathematical induction,
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Example Theorem: Everybody is rich.

Proof Consider $S(n)$: Given any collection of n persons, if at least one of them is rich, then they are all rich.

We prove $S(n)$ by mathematical induction.

Induction basis: $S(1)$ is obviously true.

Induction step: Assume that $S(k)$ is true.

Consider $k+1$ persons $\{\alpha_1, \alpha_2, \dots, \alpha_k, \alpha_{k+1}\}$

Suppose α_i is rich. Then $\{\alpha_1, \alpha_2, \dots, \alpha_k\}$ is a set of k persons. Since α_i is rich, by $S(k)$, $\alpha_1, \alpha_2, \dots, \alpha_k$ are all rich.

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Also consider $\{\alpha_1, \alpha_2, \dots, \alpha_{k-1}, \alpha_{k+1}\}$ which is a set of k persons. Since α_1 is rich, by $S(k)$, $\alpha_1, \alpha_2, \dots, \alpha_{k-1}, \alpha_{k+1}$ are all rich. Therefore, $S(k)$ implies $S(k+1)$. So by mathematical induction, $S(n)$ is true for all $n \in \mathbb{N}$.



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So by mathematical induction, $S(n)$ is true for all $n \in \mathbb{N}$.

Since there certainly exists at least one rich person,
it follows that all persons are rich. \equiv

Question: What is wrong with the proof?

Answer: The induction step that $S(k)$ implies $S(k+1)$
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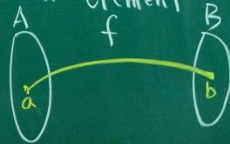
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Functions

Let A and B be nonempty sets.

Def A function $f: A \rightarrow B$ is a rule that assigns to each element $a \in A$ a unique element $b \in B$.



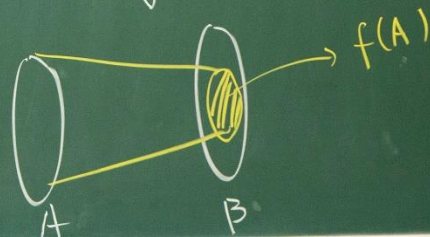
$$f(a) = b$$

A : domain of f

B : codomain of f

$$f(A) = \{f(a) : a \in A\}$$

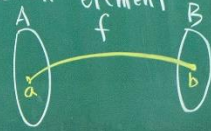
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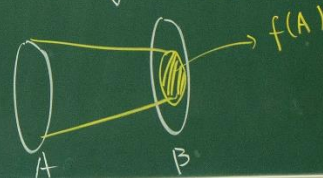
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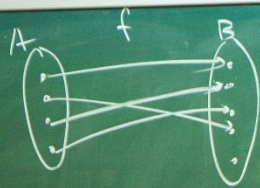
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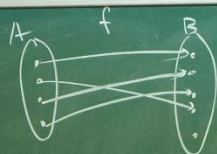
Def A function $f: A \rightarrow B$ is called **one-to-one** or **injective**, if $f(a_1) = f(a_2)$ implies $a_1 = a_2$ (or equivalently, $a_1 \neq a_2$ implies $f(a_1) \neq f(a_2)$).
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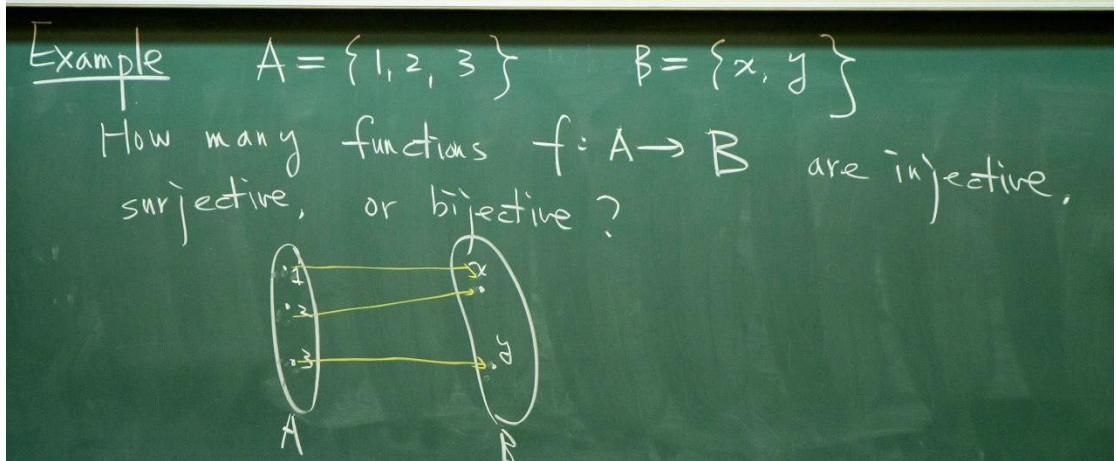


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A function $f: A \rightarrow B$ is called a **one-to-one correspondence** or **bijective** (a **bijection**) if f is both **one-to-one** and **onto** (or equivalently, both **injective** and **surjective**).



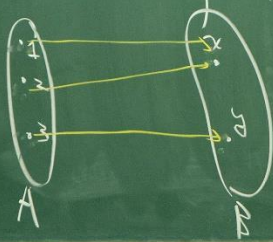
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Example $A = \{1, 2, 3\}$ $B = \{x, y\}$

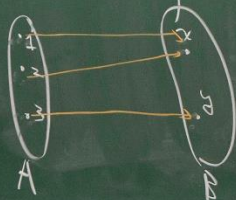
How many functions $f: A \rightarrow B$ are injective, surjective, or bijective?

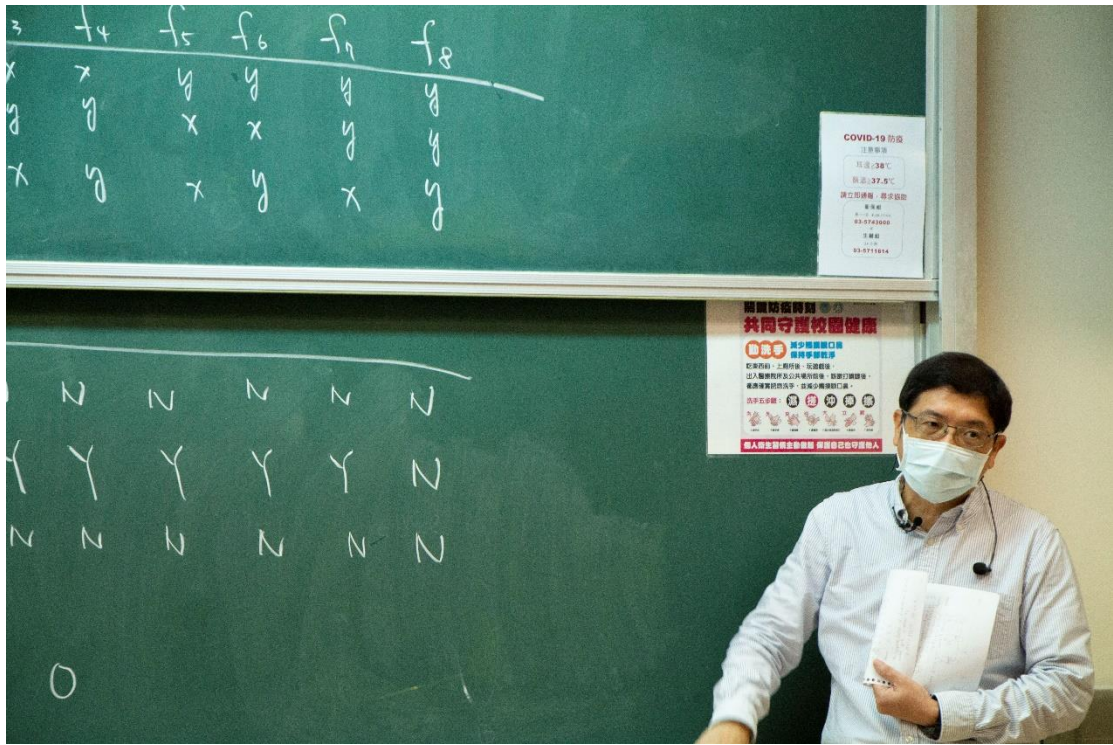


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There are totally $\underline{2 \cdot 2 \cdot 2 = 8}$ different functions $f: A \rightarrow B$.

a	f_1	f_2	f_3	f_4	f_5	f_6	f_7	f_8
1	x	x	x	x	y	y	y	y
2	x	x	y	y	x	x	y	y
3	x	y	x	y	x	y	x	x

